## Methods of solving sparse linear systems.

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## Methods of solving sparse linear systems.

Outline

- Introduction
- Sherman-Morrison formula
- Woodbury formula
- Indexed storage of sparse matrices


## Methods of solving sparse linear systems.

## Types of sparse matrixes

Band diagonal with bandwidth M


Block tridiagonal


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## Types of sparse matrixes

Doubly bordered block diagonal


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## Sherman-Morrison Formula

- The general idea of Sherman-Morrison formula is replacement of the original matrix in the sum of matrix $A$ and the product of vectors $u$ and $v$.

$$
A \longrightarrow(A+u \otimes v)
$$

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## Sherman-Morrison Formula

$$
\begin{aligned}
& (A+u \otimes v)^{-1}=\left(1+A^{-1} \cdot u \otimes v\right)^{-1} \cdot A^{-1} \\
& =\left(1-A^{-1} \cdot u \otimes v+A^{-1} \cdot u \otimes v \cdot A^{-1} \cdot u\right. \\
& \otimes v-\ldots) \cdot A^{-1} \\
& =A^{-1}-A^{-1} \cdot u \otimes v \cdot A^{-1}\left(1-\lambda+\lambda^{2}-\ldots\right) \\
& =A^{-1}-\frac{\left(A^{-1} \cdot u\right) \otimes\left(v \cdot A^{-1}\right)}{1+\lambda} \\
& \text { Where } \\
& \lambda \equiv v \cdot A^{-1} \cdot u
\end{aligned}
$$

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## Sherman-Morrison Formula

Given $\boldsymbol{A}^{\mathbf{- 1}}$ and the vectors $\boldsymbol{u}$ and $\boldsymbol{v}$, we need only to perform matrix multiplication and a vector dot product,

$$
\begin{gathered}
z \equiv A^{-1} \cdot u \\
\lambda \equiv v \cdot z
\end{gathered}
$$

to find inverse matrix

$$
A^{-1} \rightarrow A^{-1}-\frac{z \otimes w}{1+\lambda}
$$

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## Sherman-Morrison Formula

The Sherman-Morrison formula can be directly applied to a class of sparse problems. If you already have a fast way of calculating the inverse of $\boldsymbol{A}$, then this method allows you to build up a method for more complicated matrices, adding for example a row or a column at a time. Notice that you can apply the Sherman-Morrison formula more than once successively, using at each stage the most recent update of $A^{-1}$. Of course, if you have to modify every row, then you are back to an $N^{3}$ method.

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## Sherman-Morrison Formula

For some other sparse problems, the Sherman-Morrison formula cannot be directly applied for the simple reason that storage of the whole inverse matrix $A^{-1}$ is not feasible. If you want to add only a single correction of the form $\boldsymbol{u} \otimes \boldsymbol{v}$, and solve the linear system

$$
(A+u \otimes v) \cdot x=b
$$

Then you proceed as follows. Using the fast method that is presumed available for the matrix , solve the two auxiliary problems

$$
A \cdot y=b \quad A \cdot z=u
$$

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## Sherman-Morrison Formula

In terms of these solutions,

$$
x=y-\left(\frac{v \cdot y}{1+(v \cdot z)}\right) z
$$

## Methods of solving sparse linear systems.

## Sherman-Morrison Formula

## Example

Cyclic tridiagonal system:

$$
\left[\begin{array}{ccccccc}
b_{0} & c_{0} & 0 & \ldots & & & \beta \\
a_{1} & b_{1} & c_{1} & \ldots & & & \\
& & & \ldots & a_{N-2} & b_{N-2} & c_{N-2} \\
\alpha & & & \ldots & 0 & a_{N-1} & b_{N-1}
\end{array}\right] \cdot\left[\begin{array}{c}
x_{0} \\
x_{1} \\
\ldots \\
x_{N-2} \\
x_{N-1}
\end{array}\right]=\left[\begin{array}{c}
r_{0} \\
r_{1} \\
\ldots \\
r_{N-2} \\
r_{N-1}
\end{array}\right]
$$

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## Sherman-Morrison Formula

- This is a tridiagonal system, except for the matrix elements $\alpha$ and $\beta$ in the corners.
- We use the Sherman-Morrison formula, treating the system as a tridiagonal plus a correction.

$$
u=\left[\begin{array}{c}
\gamma \\
0 \\
\vdots \\
0 \\
\alpha
\end{array}\right]
$$

$$
v=\left[\begin{array}{c}
1 \\
0 \\
\vdots \\
0 \\
\beta / \gamma
\end{array}\right]
$$

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## Sherman-Morrison Formula

- Then the matrix is the tridiagonal part of initial matrix, with two terms modified:

$$
b_{0}^{\prime}=b_{0}-\gamma, \quad b_{N-1}^{\prime}=b_{N-1}-\frac{\alpha \beta}{\gamma}
$$

We solve equations

$$
A \cdot y=b \quad A \cdot z=u
$$

With the standard tridiagonal algorithm, for example sweep method.

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## Sherman-Morrison Formula

And get the solution by the formula

$$
x=y-\left(\frac{v \cdot y}{1+(v \cdot z)}\right) z
$$

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## Sherman-Morrison Formula

Acceleration with respect to Gauss method


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## Woodbury Formula

If you want to add more than a single correction term, then you cannot use Sharman-Morrison formula repeatedly, since without storing a new $A^{-1}$ you are not able to solve the auxiliary problems efficiently after the first step. Instead you need the Woodbury Formula

$$
\left(A+U \cdot V^{T}\right)^{-1}=A^{-1}-\left[A^{-1} \cdot U \cdot\left(1+V^{T} \cdot A^{-1} \cdot U\right)^{-1} \cdot V^{T} \cdot A^{-1}\right]
$$

Here $A$ is, as usual, an $N \times N$ matrix, while $U$ and $V$ are matrices with $P<N$ and usually $P \ll N$.

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## Woodbury Formula

The relation between the Woodbury formula and successive applications of the Sherman-Morrison formula is now clarified by noting that, if $U$ is the matrix formed by columns with $P$ vectors $u_{0}, \ldots, u_{P-1}$, and $V$ is the matrix formed by columns with $P$ vectors $v_{0}, \ldots, v_{P-1}$

$$
U \equiv\left[u_{0}\right] \ldots\left[u_{P-1}\right] \quad V \equiv\left[v_{0}\right] \ldots\left[v_{P-1}\right]
$$

## Methods of solving sparse linear systems.

## Woodbury Formula

Then two ways of expressing the same correction to are possible

$$
\left(A+\sum_{k=0}^{P-1} u_{k} \otimes v_{k}\right)=\left(A+U \cdot V^{T}\right)
$$

Note that the subscripts on $\mathbf{u}$ and $\mathbf{v}$ do not denote components, but rather distinguish the different column vectors.

## Methods of solving sparse linear systems.

## Woodbury Formula

Last equation reveals that, if you have $A^{\mathbf{- 1}}$ in storage, then you can either make the $P$ corrections at once by using Woodbury formula, inverting a $P \times P$ matrix, or else make them by applying ShermanMorrison formula $P$ successive times.

If you don't have storage for $A^{-1}$, then you must use Woodbury formula in the following way:To solve the linear equation

$$
\left(A+\sum_{k=0}^{P-1} u_{k} \otimes v_{k}\right) \cdot x=b
$$

## Methods of solving sparse linear systems.

## Woodbury Formula

first solve the auxiliary problems

$$
\begin{gathered}
A \cdot z_{0}=u_{0} \\
A \cdot z_{1}=u_{1} \\
\ldots \\
A \cdot z_{P-1}=u_{P-1}
\end{gathered}
$$

and construct the matrix by columns of the obtained solutions

$$
Z \equiv\left[z_{0}\right]\left[\ldots\left[z_{P-1}\right]\right.
$$

## Methods of solving sparse linear systems.

## Woodbury Formula

Next, do the $P \times P$ matrix inversion

$$
H \equiv\left(1+V^{T} \cdot Z\right)^{-1}
$$

Finally, solve one more auxiliary problem

$$
A \cdot y=b
$$

In terms of these quantities, the solution is given by the formula

$$
x=y-Z \cdot\left[H \cdot\left(V^{T} \cdot y\right)\right]
$$

## Methods of solving sparse linear systems.

## Indexed Storage of Sparse Matrices

- The first N locations of sa store diagonal matrix elements of $\mathbf{A}$, in their order. (Note that diagonal elements are stored even if they are zeros; this is at most a slight storage inefficiency, since diagonal elements are nonzero in most realistic applications)
- Each of the first N locations of ija stores the index of the array sa that contains the first off-diagonal element of the corresponding row of the matrix. (If there are no off-diagonal elements for that row, it is one greater than the index in sa of the most recently stored element of a previous row)


## Methods of solving sparse linear systems.

## Indexed Storage of Sparse Matrices

- Location 0 of ija is always equal to $\mathrm{N}+\mathrm{I}$. (It can be read (или used) to determine N )
- Location N of ija is one greater than the index in sa of the last offdiagonal element of the last row. ( It ca be read to determine the number of nonzero elements in the matrix, or the number of elements in the arrays sa and ija.) Location N of sa is not used and can be set arbitrarily.
- Entries in sa at locations $>=N+I$ contain off-diagonal elements of A ordered by rows and, within each row, ordered by columns.
- Entries in ija at locations $>=\mathrm{N}+\mathrm{I}$ contain the column number (или index) of the corresponding element in sa.


## Methods of solving sparse linear systems.

## Indexed Storage of Sparse Matrices

$$
A=\left[\begin{array}{lllll}
3 & 0 & 1 & 0 & 0 \\
0 & 4 & 0 & 0 & 0 \\
0 & 7 & 5 & 9 & 0 \\
0 & 0 & 0 & 0 & 2 \\
0 & 0 & 0 & 6 & 5
\end{array}\right]
$$

| Index k | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| ija[k] | 6 | 7 | 7 | 9 | 10 | 11 | 2 | 1 | 3 | 4 | 2 |
| sa[k] | 3 | 4 | 5 | 0 | 5 | $x$ | 1 | 7 | 9 | 2 | 6 |

## Thank you

