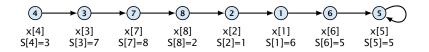
Input: a linked list given by successor pointers; a value x[i] for every list element; an operator *;

Output: for every list position ℓ the sum (w.r.t. *) of elements after ℓ in the list (including ℓ)





4.2 Parallel Prefix

Algori	thm 7 ParallelPrefix
1: for	$1 \le i \le n$ pardo
2:	$P[i] \leftarrow S[i]$
3:	while $S[i] \neq S[S[i]]$ do
4:	$x[i] \leftarrow x[i] * x[S[i]]$
5:	$S[i] \leftarrow S[S[i]]$
6:	if $P[i] \neq i$ then $x[i] \leftarrow x[i] * x[S(i)]$

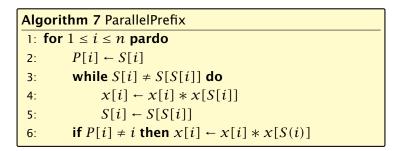
The algorithm runs in time $O(\log n)$.

It has work requirement $O(n \log n)$. non-optimal

This technique is also known as pointer jumping



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