team	wins	losses	remaining games			
i	w_i	ℓ_i	Atl	Phi	NY	Mon
Atlanta	83	71	-	1	6	1
Philadelphia	80	79	1	-	0	2
New York	78	78	6	0	-	0
Montreal	77	82	1	2	0	-

Which team can end the season with most wins?

- Montreal is eliminated, since even after winning all remaining games there are only 80 wins.
- But also Philadelphia is eliminated. Why?

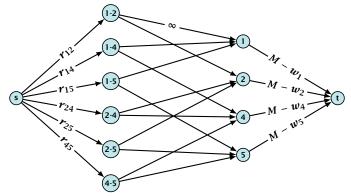


Formal definition of the problem:

- Given a set S of teams, and one specific team $z \in S$.
- Team x has already won w_x games.
- Team x still has to play team y, r_{xy} times.
- Does team z still have a chance to finish with the most number of wins.



Flow networks for z = 3. *M* is number of wins Team 3 can still obtain.



Idea. Distribute the results of remaining games in such a way that no team gets too many wins.



14.2 Baseball Elimination

Certificate of Elimination

Let $T \subseteq S$ be a subset of teams. Define

$$w(T) := \sum_{i \in T} w_i, \qquad r(T) := \sum_{i,j \in T, i < j} r_{ij}$$

wins of
teams in T remaining games
among teams in T

If $\frac{w(T)+r(T)}{|T|} > M$ then one of the teams in T will have more than M wins in the end. A team that can win at most M games is therefore eliminated.



14.2 Baseball Elimination

Theorem 83

A team z is eliminated if and only if the flow network for z does not allow a flow of value $\sum_{ij \in S \setminus \{z\}, i < j} r_{ij}$.

Proof (⇐)

- Consider the mincut A in the flow network. Let T be the set of team-nodes in A.
- If for a node x-y not both team nodes x and y are in T, then x-y ∉ A as otw. the cut would cut an infinite capacity edge.
- We don't find a flow that saturates all source edges:

$$\begin{aligned} r(S \setminus \{z\}) &> \operatorname{cap}(S, V \setminus S) \\ &\geq \sum_{i < j: \, i \notin T \lor j \notin T} r_{ij} + \sum_{i \in T} (M - w_i) \\ &\geq r(S \setminus \{z\}) - r(T) + |T|M - w(T) \end{aligned}$$

► This gives M < (w(T) + r(T))/|T|, i.e., z is eliminated.

V

Proof (⇒)

- Suppose we have a flow that saturates all source edges.
- We can assume that this flow is integral.
- For every pairing x-y it defines how many games team x and team y should win.
- The flow leaving the team-node x can be interpreted as the additional number of wins that team x will obtain.
- This is less than $M w_{\chi}$ because of capacity constraints.
- ► Hence, we found a set of results for the remaining games, such that no team obtains more than *M* wins in total.
- Hence, team *z* is not eliminated.

