## 6.3 The Characteristic Polynomial

Consider the recurrence relation:

$$c_0T(n) + c_1T(n-1) + c_2T(n-2) + \cdots + c_kT(n-k) = f(n)$$

This is the general form of a linear recurrence relation of order k with constant coefficients ( $c_0, c_k \neq 0$ ).

- T(n) only depends on the k preceding values. This means the recurrence relation is of order k.
- The recurrence is linear as there are no products of T[n]'s.
- ► If f(n) = 0 then the recurrence relation becomes a linear, homogenous recurrence relation of order k.

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# The Homogenous Case

The solution space

 $S = \{T = T[0], T[1], T[2], \dots \mid T \text{ fulfills recurrence relation} \}$ 

is a vector space. This means that if  $T_1, T_2 \in S$ , then also  $\alpha T_1 + \beta T_2 \in S$ , for arbitrary constants  $\alpha, \beta$ .

#### How do we find a non-trivial solution?

We guess that the solution is of the form  $\lambda^n$ ,  $\lambda \neq 0$ , and see what happens. In order for this guess to fulfill the recurrence we need

$$c_0\lambda^n + c_1\lambda^{n-1} + c_2 \cdot \lambda^{n-2} + \cdots + c_k \cdot \lambda^{n-k} = 0$$

for all  $n \ge k$ .

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# 6.3 The Characteristic Polynomial

#### **Observations:**

- ► The solution T[0], T[1], T[2],... is completely determined by a set of boundary conditions that specify values for T[0],...,T[k-1].
- In fact, any k consecutive values completely determine the solution.
- k non-concecutive values might not be an appropriate set of boundary conditions (depends on the problem).

#### Approach:

- First determine all solutions that satisfy recurrence relation.
- Then pick the right one by analyzing boundary conditions.
- First consider the homogenous case.

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# The Homogenous Case

Dividing by  $\lambda^{n-k}$  gives that all these constraints are identical to

$$\underbrace{c_0\lambda^k + c_1\lambda^{k-1} + c_2 \cdot \lambda^{k-2} + \dots + c_k}_{\text{characteristic polynomial } P[\lambda]} = 0$$

This means that if  $\lambda_i$  is a root (Nullstelle) of  $P[\lambda]$  then  $T[n] = \lambda_i^n$  is a solution to the recurrence relation.

Let  $\lambda_1, \ldots, \lambda_k$  be the k (complex) roots of  $P[\lambda]$ . Then, because of the vector space property

$$\alpha_1\lambda_1^n + \alpha_2\lambda_2^n + \cdots + \alpha_k\lambda_k^n$$

is a solution for arbitrary values  $\alpha_i$ .

## The Homogenous Case

#### Lemma 5

Assume that the characteristic polynomial has k distinct roots  $\lambda_1, \ldots, \lambda_k$ . Then all solutions to the recurrence relation are of the form

$$\alpha_1\lambda_1^n + \alpha_2\lambda_2^n + \cdots + \alpha_k\lambda_k^n$$
.

#### Proof.

There is one solution for every possible choice of boundary conditions for  $T[1], \ldots, T[k]$ .

We show that the above set of solutions contains one solution for every choice of boundary conditions.

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## **The Homogenous Case**

## Proof (cont.).

Suppose I am given boundary conditions T[i] and I want to see whether I can choose the  $\alpha'_i s$  such that these conditions are met:

$$\begin{pmatrix} \lambda_1 & \lambda_2 & \cdots & \lambda_k \\ \lambda_1^2 & \lambda_2^2 & \cdots & \lambda_k^2 \\ & \vdots & & \\ \lambda_1^k & \lambda_2^k & \cdots & \lambda_k^k \end{pmatrix} \begin{pmatrix} \alpha_1 \\ \alpha_2 \\ \vdots \\ \alpha_k \end{pmatrix} = \begin{pmatrix} T[1] \\ T[2] \\ \vdots \\ T[k] \end{pmatrix}$$

We show that the column vectors are linearly independent. Then the above equation has a solution.

# **The Homogenous Case**

### Proof (cont.).

Suppose I am given boundary conditions T[i] and I want to see whether I can choose the  $\alpha'_i s$  such that these conditions are met:

$$\alpha_{1} \cdot \lambda_{1} + \alpha_{2} \cdot \lambda_{2} + \cdots + \alpha_{k} \cdot \lambda_{k} = T[1]$$

$$\alpha_{1} \cdot \lambda_{1}^{2} + \alpha_{2} \cdot \lambda_{2}^{2} + \cdots + \alpha_{k} \cdot \lambda_{k}^{2} = T[2]$$

$$\vdots$$

$$\alpha_{1} \cdot \lambda_{1}^{k} + \alpha_{2} \cdot \lambda_{2}^{k} + \cdots + \alpha_{k} \cdot \lambda_{k}^{k} = T[k]$$
6.3 The Characteristic Polynomial
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# **The Homogenous Case**

#### Proof (cont.).

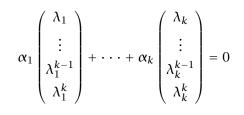
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This we show by induction:

**base case** (k = 1):

A vector  $(\lambda_i)$ ,  $\lambda_i \neq 0$  is linearly independent.

• induction step  $(k \rightarrow k + 1)$ : assume for contradiction that there exist  $\alpha_i$ 's with

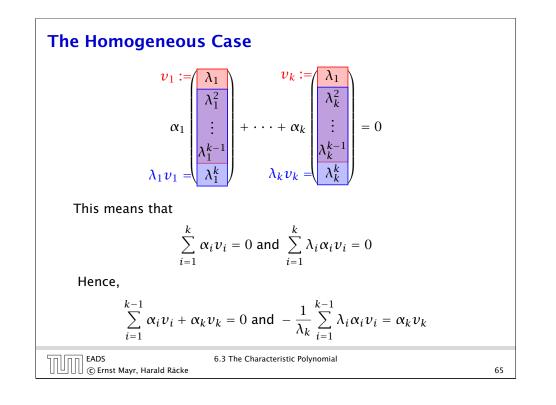


and not all  $\alpha_i = 0$ . Then all  $\alpha_i \neq 0$ !

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## The Homogeneous Case

#### What happens if the roots are not all distinct?

Suppose we have a root  $\lambda_i$  with multiplicity (Vielfachheit) at least 2. Then not only is  $\lambda_i^n$  a solution to the recurrence but also  $n\lambda_i^n$ .

To see this consider the polynomial

$$P(\lambda)\lambda^{n-k} = c_0\lambda^n + c_1\lambda^{n-1} + c_2\lambda^{n-2} + \dots + c_k\lambda^{n-k}$$

Since  $\lambda_i$  is a root we can write this as  $Q(\lambda)(\lambda - \lambda_i)^2$ . Calculating the derivative gives a polynomial that still has root  $\lambda_i$ .

This means

$$c_0 n \lambda_i^{n-1} + c_1 (n-1) \lambda_i^{n-2} + \cdots + c_k (n-k) \lambda_i^{n-k-1} = 0$$

Hence,

$$c_0 n \lambda_i^n + c_1 (n-1) \lambda_i^{n-1} + \dots + c_k (n-k) \lambda_i^{n-k} = 0$$

$$T[n] \qquad T[n-1] \qquad T[n-k]$$

# The Homogeneous Case

This gives that

$$\sum_{i=1}^{k-1} (1 - \frac{\lambda_i}{\lambda_k}) \alpha_i v_i = 0 .$$

This is a contradiction as the  $v_i$ 's are linearly independent because of induction hypothesis.

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# The Homogeneous Case

Suppose  $\lambda_i$  has multiplicity *j*. We know that

$$c_0 n \lambda_i^n + c_1 (n-1) \lambda_i^{n-1} + \dots + c_k (n-k) \lambda_i^{n-k} = 0$$

(after taking the derivative; multiplying with  $\lambda$ ; plugging in  $\lambda_i$ )

Doing this again gives

$$c_0 n^2 \lambda_i^n + c_1 (n-1)^2 \lambda_i^{n-1} + \dots + c_k (n-k)^2 \lambda_i^{n-k} = 0$$

We can continue j - 1 times.

Hence,  $n^{\ell}\lambda_i^n$  is a solution for  $\ell \in 0, \ldots, j-1$ .

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## **The Homogeneous Case**

#### Lemma 6

Let  $P[\lambda]$  denote the characteristic polynomial to the recurrence

$$c_0T[n] + c_1T[n-1] + \cdots + c_kT[n-k] = 0$$

Let  $\lambda_i$ , i = 1, ..., m be the (complex) roots of  $P[\lambda]$  with multiplicities  $\ell_i$ . Then the general solution to the recurrence is given by

$$T[n] = \sum_{i=1}^{m} \sum_{j=0}^{\ell_i - 1} \alpha_{ij} \cdot (n^j \lambda_i^n) .$$

The full proof is omitted. We have only shown that any choice of  $\alpha_{ii}$ 's is a solution to the recurrence.

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## **Example: Fibonacci Sequence**

Hence, the solution is of the form

$$\alpha\left(\frac{1+\sqrt{5}}{2}\right)^n + \beta\left(\frac{1-\sqrt{5}}{2}\right)$$

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T[0] = 0 gives  $\alpha + \beta = 0$ .

T[1] = 1 gives

$$\alpha\left(\frac{1+\sqrt{5}}{2}\right)+\beta\left(\frac{1-\sqrt{5}}{2}\right)=1 \Longrightarrow \alpha-\beta=\frac{2}{\sqrt{5}}$$

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#### **Example: Fibonacci Sequence**

T[0] = 0 T[1] = 1 $T[n] = T[n-1] + T[n-2] \text{ for } n \ge 2$ 

The characteristic polynomial is

 $\lambda^2 - \lambda - 1$ 

Finding the roots, gives

$$\lambda_{1/2} = \frac{1}{2} \pm \sqrt{\frac{1}{4} + 1} = \frac{1}{2} \left( 1 \pm \sqrt{5} \right)$$

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# Example: Fibonacci Sequence Hence, the solution is $\frac{1}{\sqrt{5}} \left[ \left( \frac{1 + \sqrt{5}}{2} \right)^n - \left( \frac{1 - \sqrt{5}}{2} \right)^n \right]$ $\frac{1}{\sqrt{5}} \left[ \left( \frac{3}{2} \right)^n - \left( \frac{1 - \sqrt{5}}{2} \right)^n \right]$

## The Inhomogeneous Case

Consider the recurrence relation:

 $c_0T(n) + c_1T(n-1) + c_2T(n-2) + \cdots + c_kT(n-k) = f(n)$ 

with  $f(n) \neq 0$ .

While we have a fairly general technique for solving homogeneous, linear recurrence relations the inhomogeneous case is different.

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## The Inhomogeneous Case

Example:

T[n] = T[n-1] + 1 T[0] = 1

Then,

T[n-1] = T[n-2] + 1  $(n \ge 2)$ 

Subtracting the first from the second equation gives,

 $T[n] - T[n-1] = T[n-1] - T[n-2] \qquad (n \ge 2)$ 

or

 $T[n] = 2T[n-1] - T[n-2] \qquad (n \ge 2)$ 

I get a completely determined recurrence if I add T[0] = 1 and T[1] = 2.

## The Inhomogeneous Case

The general solution of the recurrence relation is

$$T(n) = T_h(n) + T_p(n) ,$$

where  $T_h$  is any solution to the homogeneous equation, and  $T_p$  is one particular solution to the inhomogeneous equation.

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There is no general method to find a particular solution.

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The Inhomogeneous Case	
Example: Characteristic polynomial:	
$\underbrace{\frac{\lambda^2 - 2\lambda + 1}{(\lambda - 1)^2}} = 0$	
Then the solution is of the form	
$T[n] = \alpha 1^n + \beta n 1^n = \alpha + \beta n$	
$T[0] = 1$ gives $\alpha = 1$ .	
$T[1] = 2$ gives $1 + \beta = 2 \Longrightarrow \beta = 1$ .	
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## The Inhomogeneous Case

If f(n) is a polynomial of degree r this method can be applied r + 1 times to obtain a homogeneous equation:

 $T[n] = T[n-1] + n^2$ 

Shift:

$$T[n-1] = T[n-2] + (n-1)^2 = T[n-2] + n^2 - 2n + 1$$

Difference:

$$T[n] - T[n-1] = T[n-1] - T[n-2] + 2n - 1$$

$$T[n] = 2T[n-1] - T[n-2] + 2n - 1$$

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## **6.4 Generating Functions**

#### Definition 7 (Generating Function)

Let  $(a_n)_{n\geq 0}$  be a sequence. The corresponding

generating function (Erzeugendenfunktion) is

$$F(z) := \sum_{n=0}^{\infty} a_n z^n;$$

 exponential generating function (exponentielle Erzeugendenfunktion) is

$$F(z) = \sum_{n \ge 0} \frac{a_n}{n!} z^n.$$

T[n] = 2T[n-1] - T[n-2] + 2n - 1

Shift:

$$T[n-1] = 2T[n-2] - T[n-3] + 2(n-1) - 1$$
$$= 2T[n-2] - T[n-3] + 2n - 3$$

Difference:

$$T[n] - T[n-1] = 2T[n-1] - T[n-2] + 2n - 1$$
$$- 2T[n-2] + T[n-3] - 2n + 3$$

$$T[n] = 3T[n-1] - 3T[n-2] + T[n-3] + 2$$

and so on...

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# **6.4 Generating Functions**

#### Example 8

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1. The generating function of the sequence (1, 0, 0, ...) is

F(z)=1.

2. The generating function of the sequence (1, 1, 1, ...) is

$$F(z) = \frac{1}{1-z}.$$

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